# Dimension theory of representations of real (and complex) numbers 

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## $b$-adic representation

- Consider the b-adic representation of a real numbers $x \in[0,1]$ :

- Choosing digits from $A \subseteq\{0, \ldots, b-1\}$ we define

$$
\mathcal{D}_{\text {b-adic }}[A]:=\left\{x \in[0,1] \mid d_{i}(x) \in A\right\} .
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- If $2 \geq|A|<b$ the set $\mathcal{D}_{\text {b-adic }}[A]$ is uncountable and compact but of length zero and totally disconnected.


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## Hausdorff dimension

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- The is a natural generalization of the $n$-dimensional Lebesgue measure to non-integer dimensions, $\mathfrak{L}^{n}=c_{n} \mathfrak{H}^{n}$.
- The Hausdorff dimension is given by

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## Hausdorff (1919)

## Theorem

$$
\operatorname{dim} \mathcal{D}_{b-a d i c}[A]=\frac{\log |A|}{\log b}
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- For the upper bound just cover the set by by $|A|^{n}$ intervals of length $b^{n}$.
- For the lower bound define a probability measure by $\mu\left(l_{a_{1} a_{2} \ldots a_{n}}\right)=|A|^{-n}$. We have $x \in \mathcal{D}: \mu\left(B_{r}(x)\right) \leq c r^{\log |A| / \log b}$.

By the mass distribution principle $\mathfrak{H}^{\log |A| / \log b}(\mathcal{D}) \geq 1 / c$.

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## Prescribed frequencies in b-adic representation

- Let $\mathbf{p}=\left(p_{j}\right)$ be a probability vector on $\{0, \ldots b-1\}$. The entropy of $\mathbf{p}$ is

- Consider the set of real numbers in $[0,1]$ with given frequency of digits in the b-adic representation

- $\mathcal{F}_{\text {b-adic }}[(1 / b)]$ is the set of normal numbers to base $b$. Almost every number is normal.


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\operatorname{dim} \mathcal{F}_{b-a d i c}[\mathbf{p}]=\frac{H(\mathbf{p})}{\log b}(=: \theta)
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- Construct a measure $\mu\left(I_{d_{1} d_{2} \ldots d_{n}}\right)=p_{d_{1}} p_{d_{2}} \ldots p_{d_{n}}$.

- $s<\theta: \lim _{r \rightarrow \infty} \mu\left(B_{r}(x)\right) / r^{s}=0 \Rightarrow \mathfrak{H}^{s}(\mathcal{F})=\infty$
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## A modification of the dyadic representation

- Represent a real number $x \in(0,1]$ by a sequence in $\mathbb{N}^{\mathbb{N}}$ :

- $n_{i}(x)$ is the distance between two digits 1 in the dyadic expansion.
- For $A \subseteq \mathbb{N}$ consider the set of real numbers $\mathcal{D}_{\text {m.dyadic }}[A]$ with digits in $A$.
- $\mathcal{D}_{\text {m. dyadic }}[\{1,2\}]=\left\{x \mid n_{i}(x)=0 \Rightarrow n_{i+1}(x)=1\right\}$ is called the golden Markov set.


## A modification of the dyadic representation

- Represent a real number $x \in(0,1]$ by a sequence in $\mathbb{N}^{\mathbb{N}}$ :

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x=\sum_{i=1}^{\infty} 2^{-\left(n_{1}(x)+\cdots+n_{i}(x)\right)}, \quad n_{i}(x) \in \mathbb{N}
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## Theorem

The Hausdorff dimension d of $\mathcal{D}$ m.dyadic $[A]$ is $d$ is given by


- For $A=\{1, \ldots n\}: d=\log (s) / \log (2)$ where $s$ is given by the solution $s \in(1,2)$ of $s^{n}-s^{n-1} \cdots-s-1=0$.
- For the golden Markov set: $d=\log ((\sqrt{5}+1) / 2) / \log 2$.
- For $A=\{n j \mid n \in \mathbb{N}\}$ we have $d=1 / j$.
- For $A=\left\{n j+m \mid n \in \mathbb{N}_{0}\right\} d$ is given by $2^{-d j}+2^{-d m}=1$.


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Consider the set of real numbers $\mathcal{F}_{\text {m.dyadic }}[\mathbf{p}]$ with frequency of digits given by by a probability vector $\mathbf{p}$ with expectation $E(\mathbf{p})$ and entropy $H(\mathbf{p})$

Theorem


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- $\operatorname{dim} \mathcal{F}_{\text {m.dyadic }}\left[\left(1 / 2,1 / 4, \ldots, 1 / 2^{n}, \ldots\right)\right]=1$

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## Continued fraction representation

- Represent a real number $x \in(0,1)$ by a continued fraction:

- Consider the set of numbers $\mathcal{D}_{\text {con. }}[A]$ with digits in $A$.


## Theorem


for $n>8$.

## Continued fraction representation

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## Jarnik (1929):

Theorem

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1-\frac{4}{n \log 2} \leq \operatorname{dim}_{H} \mathcal{D}_{\operatorname{con} .}[\{1, \ldots, n\}] \leq 1-\frac{1}{8 n \log n}
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- The calculation $\mathcal{D}_{\text {con }}$. $A$ ] has been addressed over the years:
- Today we have an efficient algorithm du to Jenkinson Pollicott (2001). Especially:
- $\operatorname{dim} \mathcal{D}_{\text {con. }} .[\{1,2\}]=0.531280506277 \ldots$ (54 digits known)
- $\operatorname{dim} \mathcal{D}_{\text {con. }} .\{\{1,2,3\}]=0.7046 \ldots$
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## As a corollary of Jarnik's dimension estimate:

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\operatorname{dim}_{H}\left\{x \in(0,1) \mid\left(n_{k}(x)\right) \text { is bounded }\right\}=1
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## Complex continued fractions

- For $z \in \mathbb{C}$ consider the Hurwiz continued fraction

- The digits $c_{j}$ are given by

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z_{j+1}=1 / z_{j}-\left[1 / z_{j}\right]=1 / z_{j}-c_{j}
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with $c_{0}=[z]$ and $z_{0}=z-c_{0}$ where [.] denotes rounding to the nearest element of $\mathbb{Z}[i]$.

- For $A \subseteq \mathbb{N}[i]$ consider the set of Hurwitz continued fractions $\mathcal{D}_{\text {comlex }}[A]$ with digits in $A$.


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\begin{gathered}
d<\operatorname{dim} \mathcal{D}_{\text {comlex }}[A]<D \\
\sum_{a+b i \in A}\left(\frac{1}{a^{2}+b^{2}}\right)^{D}=1 \\
\sum_{a+b i \in A}\left(\frac{1}{a^{2}+b^{2}+(1+\sqrt{2}) \max \{a, b\}+1}\right)^{d}=1
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- $0.21<\operatorname{dim} \mathcal{D}_{\text {comlex }}[\{3+i, 2+4 i\}]<0.27$
- $0.49<\operatorname{dim} \mathcal{D}_{\text {comlex }}[\{2+2 i, 3+2 i, 2+3 i, 3+3 i\}]<0.61$
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## Continued logarithm representation

- Consider the continued logarithm representation to base $m \geq 3$ of $x \in[0,1]$ :
$x=\lim _{n \rightarrow \infty} \log _{m}\left(d_{n}(x)+\log _{m}\left(d_{n-1}(x)+\log _{m}\left(\cdots+\log _{m}\left(d_{1}(x)\right)\right)\right.\right.$
with digits in $\{1, \ldots, m-1\}$.
- The representation is unique up to a countable set and in almost all numbers all digits appear.
- For $m \geq 4$ choosing digits from $A \subseteq\{1, \ldots, m-1\}$ we define

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# Let $\left[\left(d_{1}, \ldots, d_{n}\right)\right](x)=\log _{m}\left(d_{n}+\log _{m}\left(d_{n-1}+\cdots+\log _{m}\left(d_{1}+x\right).\right)\right.$. 

## Theorem

$$
\begin{gathered}
L_{n} \leq \operatorname{dim}_{H} \mathcal{D}_{c . l o g}[A] \leq U_{n} \\
\text { for all } n \geq 1 \text {, where } U_{n} \text { and } O_{n} \text { are given by } \\
\sum_{d_{1}, \ldots, d_{n} \in A}\left[\left(d_{k}\right)\right]^{\prime}(1)^{U_{n}}=1 \sum_{d_{1}, \ldots, d_{n} \in A}\left[\left(d_{k}\right)\right]^{\prime}(0)^{L_{n}}=1
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For $m=4$ using Mathematica we get

$$
\begin{aligned}
& \operatorname{dim}_{H} \mathcal{D}_{C \cdot l \log }[\{1,2\}]=0.81 \pm 0.01 \\
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For an arbitrary continued logarithm expansion to base $m \geq 3$ we consider the set of real numbers $\mathcal{F}_{c . \log }[\mathbf{p}]$ with frequency of digits given by by a probability vector $\mathbf{p}$.

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for all $\mathbf{p}$ (!).
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Jörg Neunhäuserer
Dimension theory of representations of numbers

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for all $\mathbf{p}$ (!).
For $m=3$ the upper bound look as follows


For an arbitrary continued logarithm expansion to base $m \geq 3$ we consider the set of real numbers $\mathcal{F}_{c . \log }[\mathbf{p}]$ with frequency of digits given by by a probability vector $\mathbf{p}$.

## Theorem

$$
\operatorname{dim}_{H} \mathcal{F}_{c \cdot \log }[\mathbf{p}] \leq c<1
$$

for all $\mathbf{p}$ (!).
For $m=3$ the upper bound look as follows


## Thanks for Your Attention



